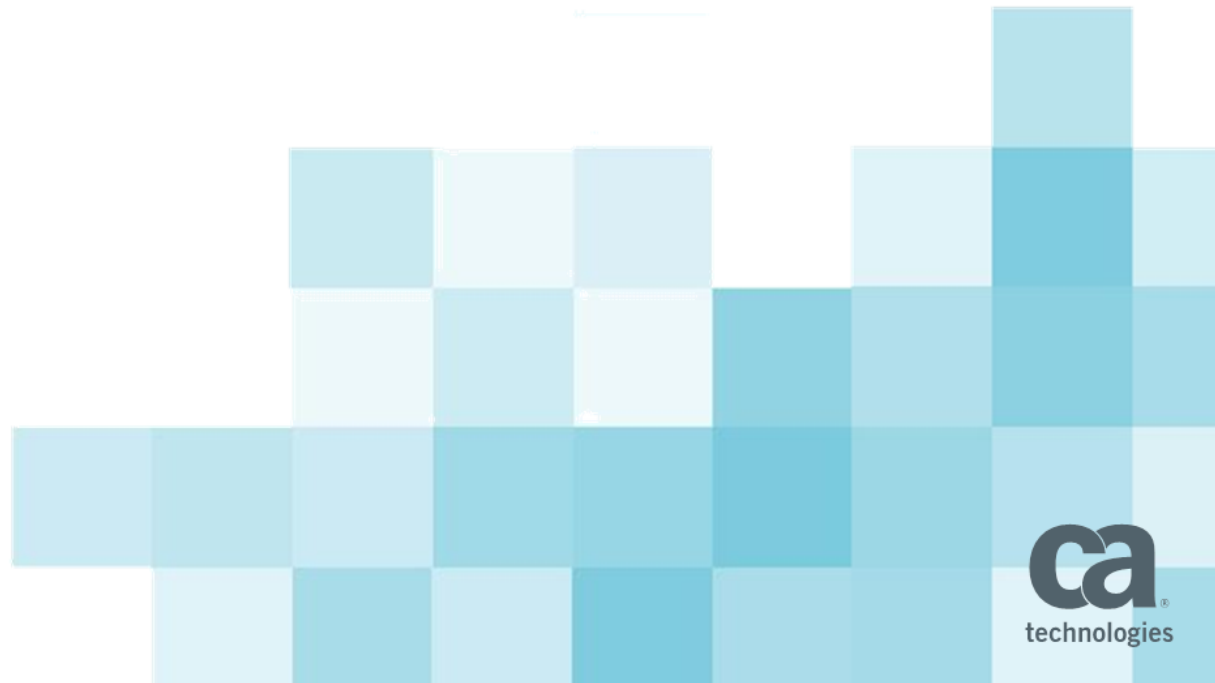


Recent Changes in DB2 for z/OS Logging

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Agenda

- What is logged?
- Reordered Row Format
- Not Logged Tablespaces
- LOBs and XML data
- LRSN Spin Avoidance
- Summary of Logging changes in DB2 9, 10, 11 and 12

What data is logged during a standard UPDATE?

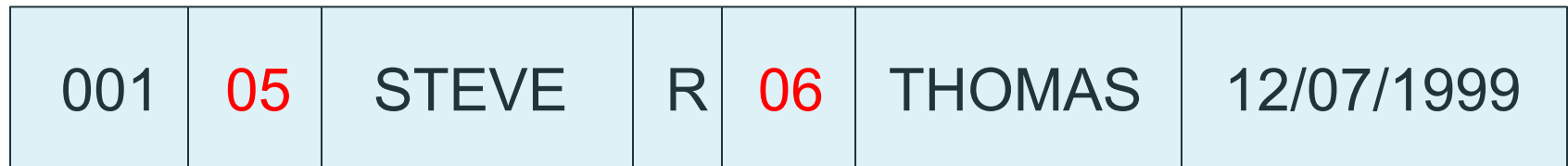
- Assuming Data Capture Changes is not used...
 - If DCC the full image of either the UNDO or REDO record is logged
 - Partial image of the other
- Fixed Data - from first to last updated byte
 - Also true for Variable data if the length does not change
- Variable Data – from first updated byte to end of row
 - Bearing in mind the first updated byte will be the column length field
- Leads to long standing good design practice
 - Locate Variable Columns at the end of the row
 - With the most updated ones towards the back end
 - Locate heavily updated fixed length columns together
- Usually only implemented when table first created
 - Rules tend to get relaxed as columns are added

Reordered Row Format (RRF)

- New Row format introduced in DB2 9 NFM
- Automatically implements some of these best practices
- Columns internally reordered on the data page
- All Fixed Columns at the beginning of the row
- All variable columns stored at the end of the row
- No lengths – uses offsets into row instead
 - Both offsets and lengths are 2 bytes so row length identical
 - Row header is not included in the offset
- Logical Column location in the table unaffected
- No Application impact – not even a Rebind required

Basic Row Format (BRF)

| | | |
|----------------|-------------|----------|
| STAFF_NUMBER | INTEGER | NOT NULL |
| FIRST_NAME | VARCHAR(30) | NOT NULL |
| MIDDLE_INITIAL | CHAR(1) | NOT NULL |
| LAST_NAME | VARCHAR(30) | NOT NULL |
| DATE_JOINED | DATE | NOT NULL |

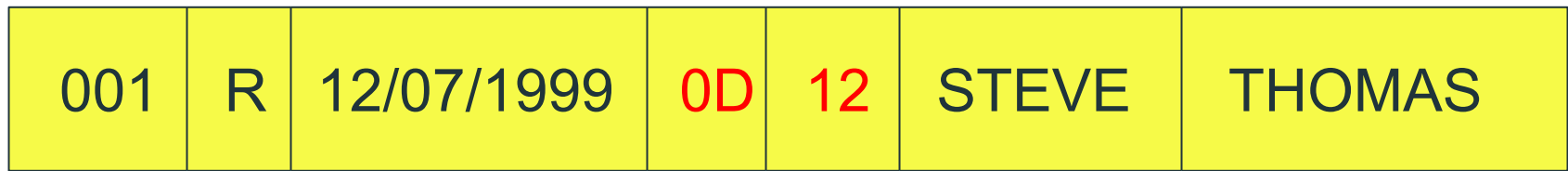


2 byte Lengths
(in hex)

Note: Integer and Date columns both use 4 bytes of internal storage

Let's look at the same row in Reordered Row Format

| | | |
|----------------|-------------|----------|
| STAFF_NUMBER | INTEGER | NOT NULL |
| FIRST_NAME | VARCHAR(30) | NOT NULL |
| MIDDLE_INITIAL | CHAR(1) | NOT NULL |
| LAST_NAME | VARCHAR(30) | NOT NULL |
| DATE_JOINED | DATE | NOT NULL |



2 byte Offsets (hex)

Note: Integer and Date columns both use 4 bytes of internal storage

RRF usually leads to reduced logging

| | | | | | | |
|-----|----|-------|---|----|--------|------------|
| 001 | 05 | STEVE | R | 06 | THOMAS | 12/07/1999 |
|-----|----|-------|---|----|--------|------------|

| | | | | | | |
|-----|----|--------|---|----|--------|------------|
| 001 | 06 | STEVEN | R | 06 | THOMAS | 12/07/1999 |
|-----|----|--------|---|----|--------|------------|



| | | | | | | |
|-----|---|------------|----|----|-------|--------|
| 001 | R | 12/07/1999 | 0D | 12 | STEVE | THOMAS |
|-----|---|------------|----|----|-------|--------|

| | | | | | | |
|-----|---|------------|----|----|--------|--------|
| 001 | R | 12/07/1999 | 0D | 13 | STEVEN | THOMAS |
|-----|---|------------|----|----|--------|--------|



When might it get worse?

- Multiple Variable columns, Update one of the later ones
 - Example – 5 columns and update only the 3rd



- Mainly impacts non-compressed data - if the data is compressed DB2 usually logs the whole row anyway
- Most customers should gain much more than they lose

Test Results verify this

- What we did...
 - Created a table with 10 columns
 - 7 were VARCHAR of various sizes
 - Loaded 1,000 rows using both BRF and RRF
 - Updated the various columns on every row
 - Measured the results using a Log Analysis tool
- In most cases log volume between flat and 50% lower
 - Exception was updating the last columns
 - Worst case example we managed to double the log volume
 - Extreme example designed to cause maximum pain
- **Conclusion: RRF should help in most cases**
 - Beware situations with multiple VARCHAR columns

Not Logged Tablespaces

- Added to provide an SQL equivalent to LOG NO utilities
 - Reduce Logging volume when data can be rebuilt
 - Never intended to be a performance option
- Prevents Logging of UNDO and REDO records
 - Control records are still logged (OPEN, CLOSE etc.)
 - So are Open URID records
 - Needed for Data Sharing and Long running URID messages
- Cannot be set for Catalog, XML or Workfiles
- Ideal applications include MQTs & Summary objects
- Default is always to Log updates
 - Except for Workfile tablespaces

Not Logged Tablespaces

- Incompatible with Data Capture Changes
- No SYSLGRNX entries created
- No SHRLEVEL CHANGE Copy or Reorg
 - Must use SHRLEVEL REFERENCE
- QUIESCE WRITE (YES)
 - Issues a Drain, Flushes the Bufferpool as normal
 - But no SYSCOPY row is created
- When you Update a NOT LOGGED object
 - Status is changed to ICOPY (non restrictive)

-CANCEL THREAD and Aborted URIDs with NOT LOGGED

- Cancelling a thread on a NOT LOGGED object
 - Leaves the object in LPL
 - And in RECP, RBDP or AUXW depending on the type
 - Consider using –CANCEL THREAD..... NOBACKOUT
 - But remember the Logged objects as well!
- LPL can only be removed manually
 - No Automatic LPL recovery
 - Recover, Refresh (MQT), Load or Drop/Recreate
 - Can also issue DELETE without WHERE or TRUNCATE
- Same consideration applies to ROLLBACK
 - Except for LOB objects – see later
- Need to ensure no Duplicate Key or RI violations

Linking Objects

- NOT LOGGED usually managed via Base Tablespace
 - This is called being **LINKED**
- XML and Indexes are always Linked
- LOB data can become Unlinked
 - Under certain circumstances
 - See next slide for a graphical example
- Look for LOG column in SYSIBM.SYSTABLESPACE
 - Base Tablespace is set to 'Y' or 'N'
 - Linked Objects are set to 'Y' or 'X'
 - LOB objects can be 'N' if they are Unlinked

Let's see an example – LOG column values

| | | | | | |
|-----|---|------------------------------|----------------------|-------------------------------|------------------------------|
| TS | Y | N | Y | Y | N |
| IX | Y | X | Y | Y | X |
| LOB | Y | X | Y | N | N |
| XML | Y | X | Y | Y | X |
| | | TS Not Logged | TS Logged | LOB Not Logged | TS Not Logged |

Potential Impact on Data Sharing

- NOT LOGGED data is not protected by the Log
 - So it needs to be externalized from Buffers quickly
- Accomplished by treating them differently
 - Read Only Check times (PCLOSEN & PCLOSET) effectively set to 1
 - DB2 treats the object like it was Read Only
 - Buffers are Forced much faster
- Good for data integrity but may have DS impact
 - Objects switch GBP Dependency frequently if updated
- Another good reason to limit using this feature to objects that are genuinely read only

Large Objects (LOBs)

- LOBs being used a lot more today, including in Catalog
- LOB data has always supported NOT LOGGED
 - Syntax changed from LOG YES/NO in DB2 9
 - Old syntax is still recognized and supported
- Old 1Gb restriction for LOB data logging lifted in DB2 9 for z/OS
- Some data is always logged
 - System Pages (Space Map) even if LOG NO or NOT LOGGED used
- For LOG YES LOBs, not everything is logged
 - Only the REDO image of data is logged
 - INSERT pages, AFTER pages of UPDATES
 - UNDO images not logged
 - DELETE pages, BEFORE pages of UPDATES

How does LOB logging differ?

- Size of LOB data is the obvious consideration
- LOB Data in an Auxiliary TS is never updated in place
 - Old data marked as deleted and new data added
 - So the old data is still available if required
 - Known as Shadow Copy Recovery
- No Logging of UNDO records for LOB data
 - ROLLBACK/BACKOUT recovery for LOB data not possible
 - SQL ROLLBACK works because the space maps are logged
- Also no Logging of LOB data for DELETES
 - Space maps provide information to allow roll forward recovery
- Remember NOT LOGGED data is Forced at Commit
 - May impact Application response times

LRSN Spin Avoidance

- Before DB2 9 each Log record required a unique LRSN
 - Could result in DB2 waiting until new LRSN available
 - Each LRSN represents around 16 microseconds
 - High CPU overhead and Log Latch Contention
- DB2 9 NFM allowed duplicate LRSNs
 - Must be successive Log records from the same member
 - Must be for different pages (commonly data and index pages)
 - Consecutive Log records for same page require unique LRSN
- DB2 10 NFM extended this capability
 - Consecutive Inserts for the same data page can share LRSN
 - Does not apply to Updates, Deletes or same index page
- Very useful for multi-row Inserts with no/few indexes
- DB2 11 Extended LRSN virtually eliminates this as a problem

DB2 9 – Major Logging related changes

- Archive Log process runs in 31 bit mode (in MSTR)
 - Reduces risk of Storage related failures when running Archive Logs
 - Allows much larger Log Buffers to be used
 - Also attempts to fill next Buffer while processing previous one
- Extended BSDS became compulsory
 - Supports 10,000 Archive and 93 Active Log pairs
- Increases in storage to improve performance
 - Active Log input Buffers increased from 15 to 120
 - Improves Fast Log Apply performance by up to 100%
 - Archive Log input buffer increased – 10 tracks per stripe
- All LOB Tablespaces can be Logged
 - Limited to <1Gb LOBs in previous releases

DB2 9 – Archive Log improvements

- Archive Log access changed from BDAM to BSAM
 - Allows features of DFSMS Extended Format datasets to be used
 - Archive Logs can now be striped
 - Compression ca also be used
 - May be a viable alternative to striping
 - See DB2 10 improvements which may also affect this decision
- Archive Logs can use Extended Volumes (EAV)
 - Specify DSNTYPE=LARGE
 - Prevents multi-volume or tape based archives
 - These were required if you were using 4GB Active Logs

DB2 10 – Latch Class 19

- Critical for data consistency
 - Used to Serialize updates

- Single latch per subsystem used before DB2 10:
 1. Latch obtained
 2. RBA range reserved
 3. Log record moved into the Log Buffer
 4. Latch released

- Now Latch is released once Buffer space reserved
 - Multiple log records can be written in parallel
 - Improves potential logging rates
 - Available in CM Mode

DB2 10 – Log Buffers

- Log Buffers are now fixed in Storage
 - DSNZPARM OUTBUFF – now defaults to 4MB (was 400KB)
- 4MB should be enough for most sites
 - IFCID 001 can indicate when your value may be too small
 - QJSTWTB – Log writes waits due to no available log buffers
 - Care required as all storage defined is now fixed
- Potential reasons to increase this include:
 - Improved ROLLBACK – hopefully not relevant!
 - CONSISTENT COPY
 - Products that use IFCID306 log reading interface
 - E.g. DB2 Data Replicator for z/OS, CONCURRENT REORGs

DB2 10 – Log I/O Enhancements

- More log writes are asynchronous and in parallel
 - In DB2 9 DB2 performs re-writes of each Log CI serially
 - Typically at COMMIT (Forced Writes)
 - Avoided risk of data loss in the event of dual log I/O failures
 - But led to more waits for Log Writes at Commit
 - DB2 10 allows these to be performed in Parallel
 - Change made possible by improvements in Disk technology
 - No loss of data integrity

- Hardware and Channel Technology improvements
 - DS8800 Disks
 - High Performance FICON (HPF) I/O protocols
 - Can be used for Log I/O > 64KB on a z196

DB2 10 I/O Performance

- Taken together these changes make a big difference
 - Redbook reports 50% reduction in Log suspend time for 3 page Synchronous Writes (e.g. at Commit)
 - Absolute difference maintained for more pages
 - The relative Percentage savings are lower
- Maximum Log throughput now reported at up to 180MB/sec
 - Compared to around 100MB/sec on a DS8300 using DB2 9
- May prevent the need to use Log Striping?

DB2 10 - Dynamically adding an Active Log

- Before DB2 10 had to use DSNJU003 – needs DB2 shutdown
- New –SET LOG NEWLOG syntax added
 - Available in CM mode - Remember to add both Active log pairs
 - This is a permanent change
 - RUN DSNJLOGF (to preformat) before adding
- Can get you out of jail in an emergency
 - Archive Log fails or hangs and you're running out of active logs
 - If Log fills and you need to shut down DB2
- Look for message DSNJ110E if archive log waits
 - -ARCHIVE LOG CANCEL OFFLOAD may be required
 - Cancels hanging log offload and starts the process again

DB2 10 - Improved Checkpoint Flexibility

- DB2 9 checkpoints using Time or Number of log records
- Neither is ideal in all circumstances
- Time is usually best choice for most customers
 - But at busiest times there may be too many log records
 - This elongates subsystems restart time
- DB2 10 can checkpoint on either value
 - Set DSNZPARM CHKTYPE=BOTH
 - And provide values for both CHKFREQ and CHKLOGR
- Setting can be changed using –SET LOG command

Other DB2 10 changes

- Must use extended format BSDS
 - Will have been done already if you're migrating from DB2 9
 - Install CLIST does this if it hasn't been done already
 - But doesn't resize the BSDS
 - If migrating from V8 then you need to check - DSNJU004
 - Look for the DSNJCNVB job message
 - Better option is to track down a V8 Install Guide

- Log Apply Storage default changed to 500MB
 - DSNZPARM LOGAPSTG – used to be 100MB
 - Introduced by PM31641 after GA date

DB2 11 – Extended RBA/LRSN

- Expanded RBA/LRSN **optionally** available
- Will increase RBA from 6 to 10 bytes
 - Basically a full word available at the front of the existing 6 byte RBA
- Will also increase LRSN from 6 to 10 bytes
 - One byte added before existing LRSN
 - IBM reckon >30,000 years availability
 - My math – ~36,352 years!
 - Three bytes precision added at the end
 - Current LRSN has about 16 microsecond precision (10^{-6})
 - Current z/OS STCKE has precision to 244 picoseconds (10^{-12})
 - Expanded LRSN to 953 femtoseconds (10^{-15})
- Should virtually eliminate LRSN Spin issues for most sites

DB2 11 – Other Log related changes

- Log Reads and Writes are now zIIP eligible
- Log Output Buffer moved from MSTR to 64-bit common area
 - Avoids a cross-memory call to MSTR when updating
 - Does not require a REBIND
- IBM Reorg writes new Compression Dictionary to Log
 - Allows products such as IBM CDC to continue without pause
- Support for Not-Logged DGTs
- Re-activation of Fast Log Apply during Restart
 - Accidentally disabled during life of DB2 9!
- Reductions in Log Writes due to Column Insert processing and Synchronous Log Writes during Index Structure modifications
- Further improvements to Log Latch Class 19

DB2 12 – Logging related changes


- 7-byte RIDs introduced by RPN Partitioning
 - Used from CM mode – needs fallback toleration PTF for DB2 11 support
- Support for Large Active Log datasets
 - Up to 768Gb – should be enough for most sites!
- Support for z/OS Hyperwrite (PPRC Log Write Accelerator)
 - Up to 30% reduction in Log Write latency using PPRC
 - Retrofitted to DB2 10 and 11
- Further significant reductions in Log Latch contention
 - IBM report up to 41% CPU and 6% Elapsed reductions
- Removed Log Write Force when caching Identity Columns and Sequences in Data Sharing



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